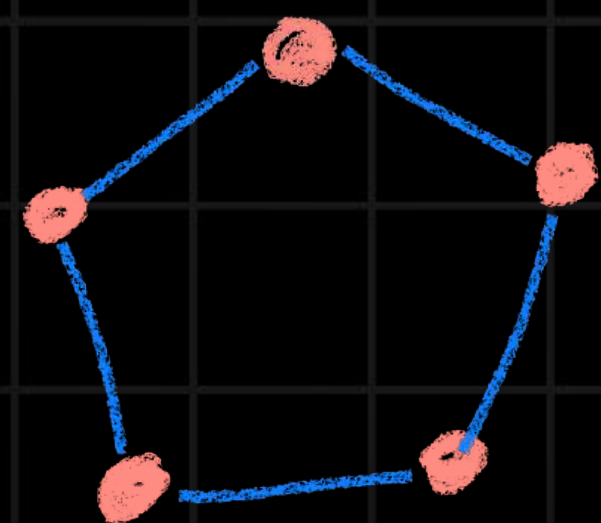


Graph Constructions as Free Random Variables

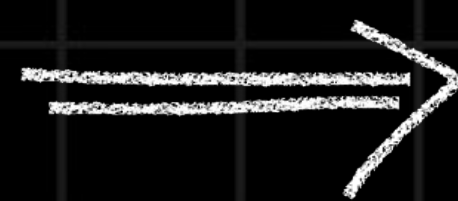
Gal Maor (TAU)

Joint works with
Gil Cohen, Itay Cohen, Yuval Peled

Spectral Expanders 101



G



$$\begin{pmatrix} \vdots & \vdots & \vdots & \vdots & \vdots \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ \vdots & \vdots & \vdots & \vdots & \vdots \end{pmatrix}$$

A_G

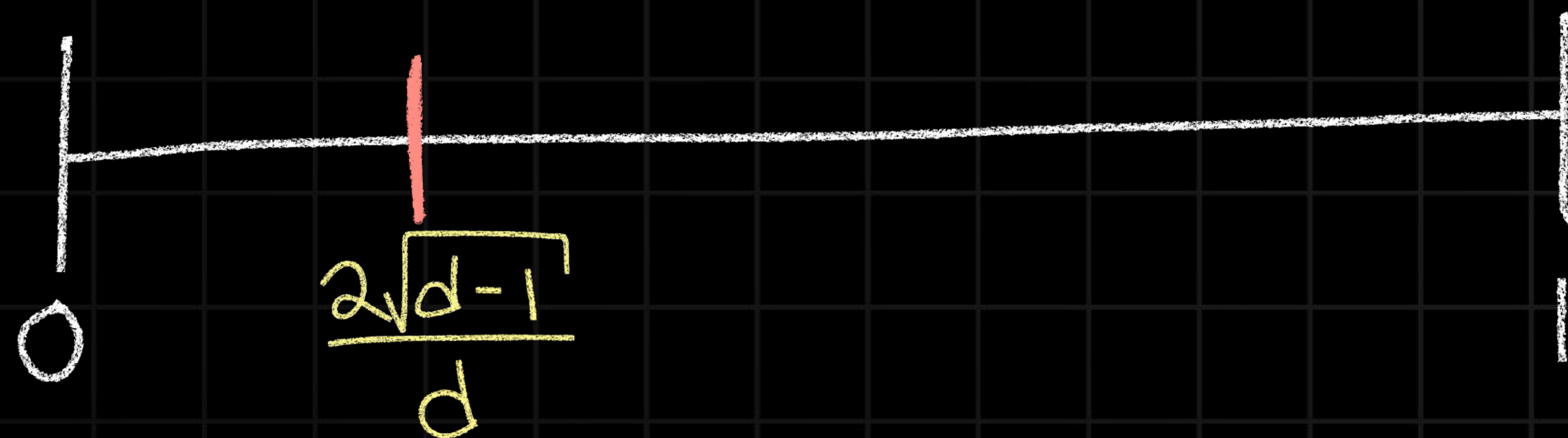
Eigenvalues of A_G :

$$d = \lambda_1 \geq \lambda_2 \geq \dots \geq \lambda_n \geq -d$$

$$\lambda(G) = \max_{i \geq 2} \{|\lambda_i|\} \in [0, d]$$

$$\omega_G = \frac{\lambda(G)}{d} \in [0, 1]$$

Alon-Boppana: $\lambda(G) \geq 2\sqrt{d-1} - o_n(1)$



$$\frac{1}{d}A_G = J + E$$

perfect
randomness

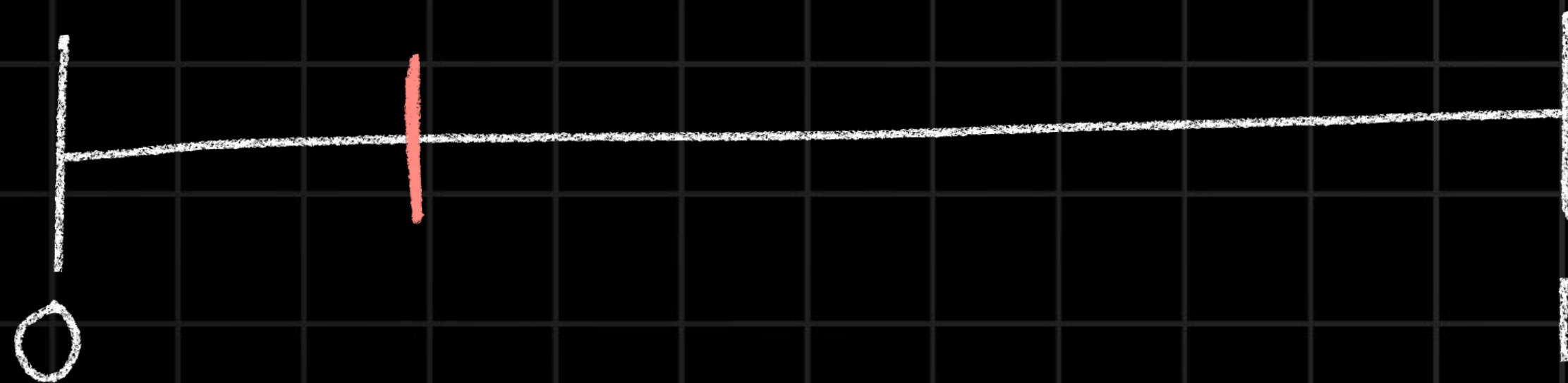
$$\|E\| = \omega_G$$

Spectral Expanders 101

Alon-Boppana: $\lambda(G) \geq 2\sqrt{d-1} - o(1)$

$$\lambda(G) = \max_{i \geq 2} \{|\lambda_i|\} \in [0, d]$$

$$w_G = \frac{\lambda(G)}{d} \in [0, 1]$$

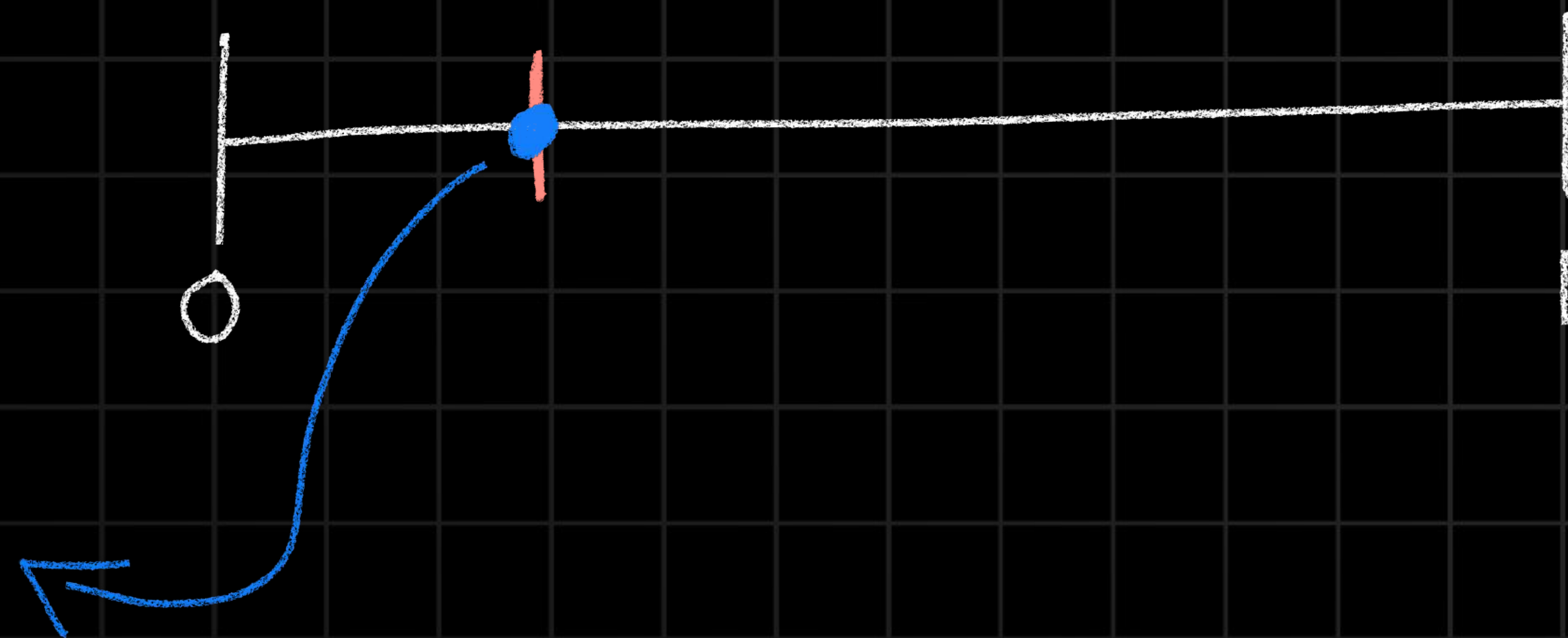


Spectral Expanders 101

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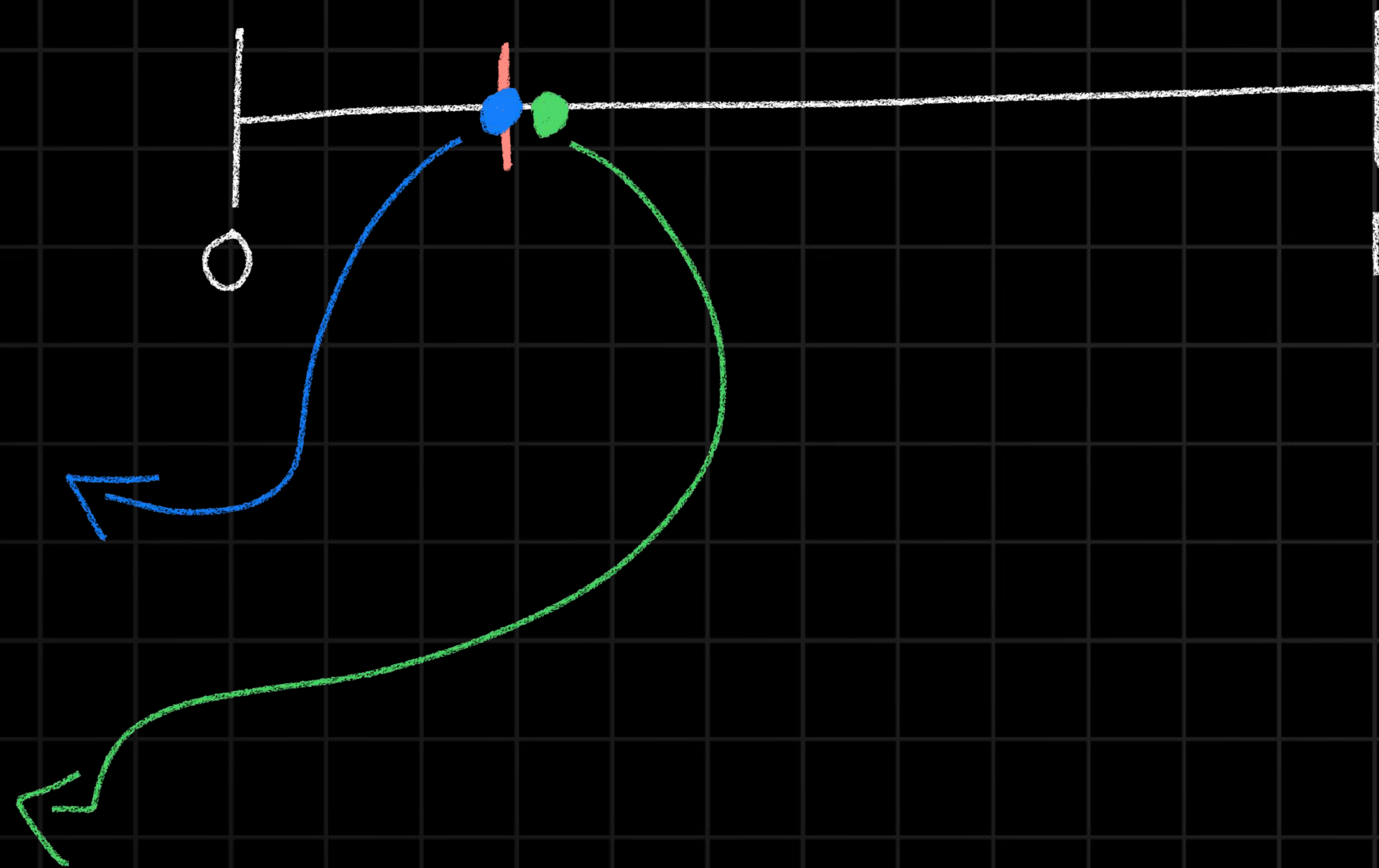
Ramanujan
[LPS'88]
[MSS'15]

Spectral Expanders 101

Alon-Boppana: $\lambda(G) \geq 2\sqrt{d-1} - o(1)$

$$\lambda(G) = \max_{i \geq 2} \{|\lambda_i|\} \in [0, d]$$

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Ramanujan

[LPS'88]

[MSS'15]

Random

[Fri'06]

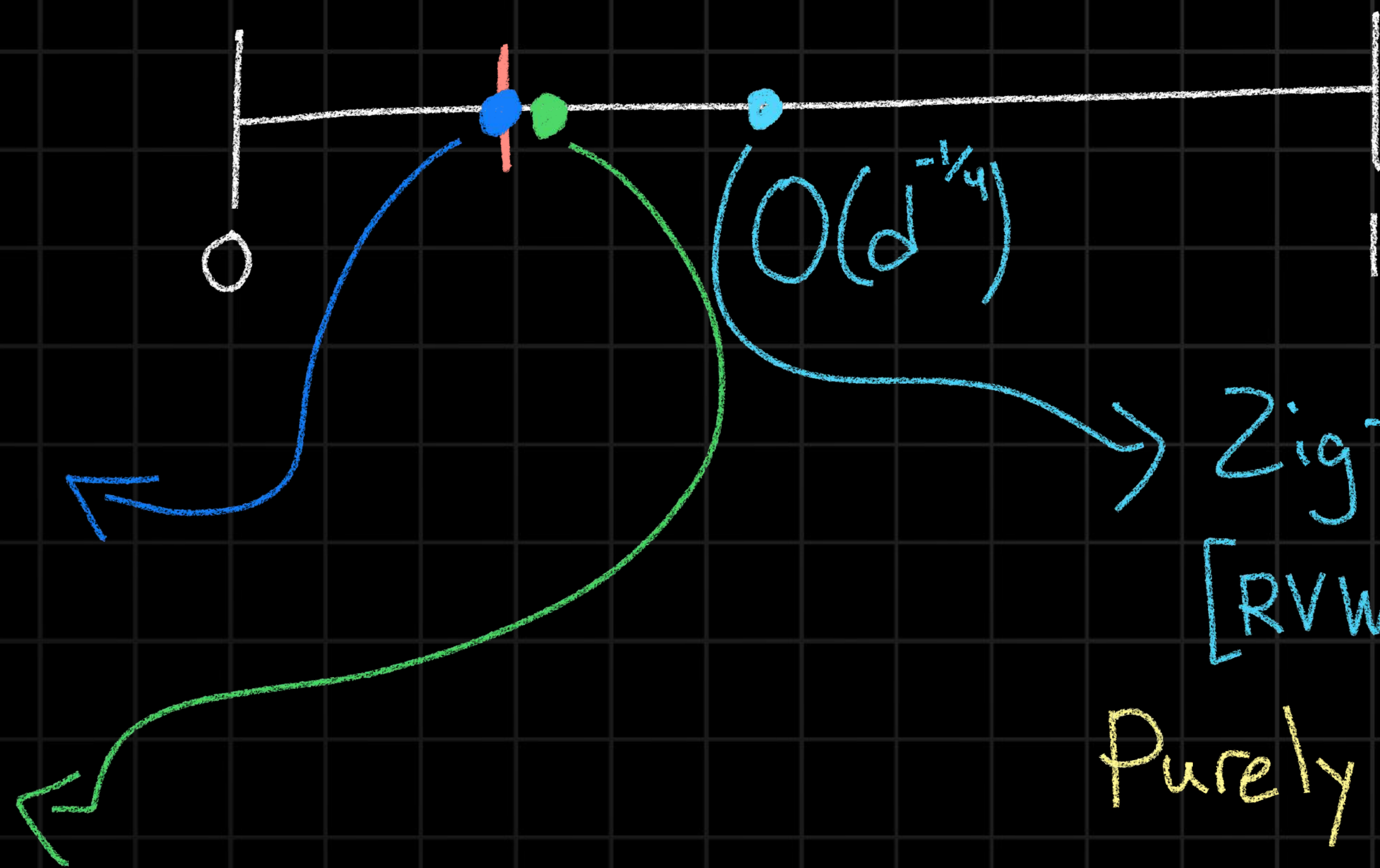
[Bor'15]

Spectral Expanders 101

Alon-Boppana: $\lambda(G) \geq 2\sqrt{d-1} - o(1)$

$$\lambda(G) = \max_{i \geq 2} \{|\lambda_i|\} \in [0, d]$$

$$\omega_G = \frac{\lambda(G)}{d} \in [0, 1]$$



Ramanujan

[LPS'88]

[MSS'15]

Random

[Fri'06]

[Bor'15]

Zig-Zag construction

[RVW'00]

Purely combinatorial!

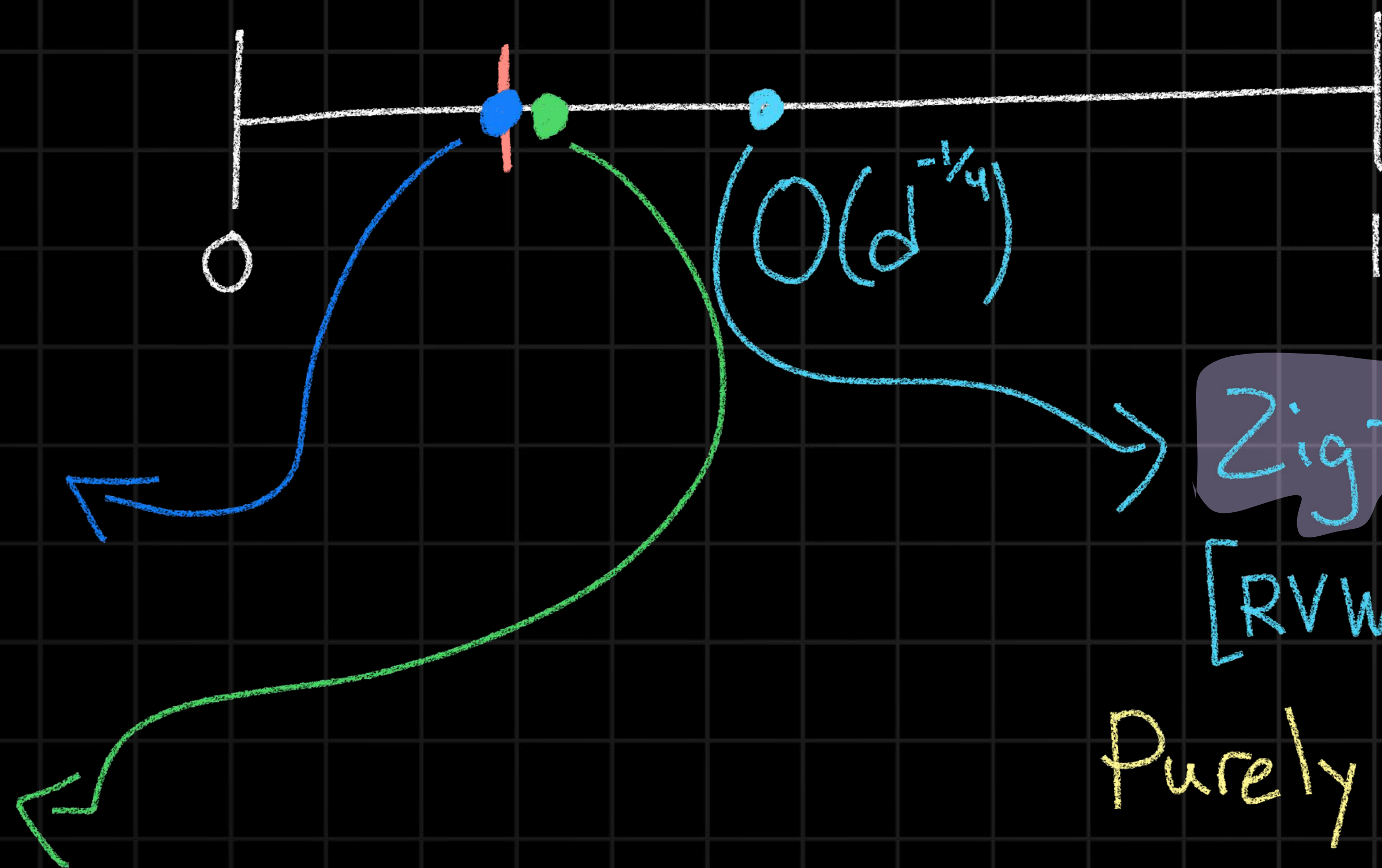
Can it be improved?

Spectral Expanders 101

Alon-Boppana: $\lambda(G) \geq 2\sqrt{d-1} - o(1)$

$$\lambda(G) = \max_{i \geq 2} \{|\lambda_i|\} \in [0, d]$$

$$\omega_G = \frac{\lambda(G)}{d} \in [0, 1]$$



Ramanujan

[LPS'88]

[MSS'15]

Random

[Fri'06]

[Bor'15]

Zig-Zag construction

[RVW'00]

Purely combinatorial!

Can it be improved?

MSS Overview

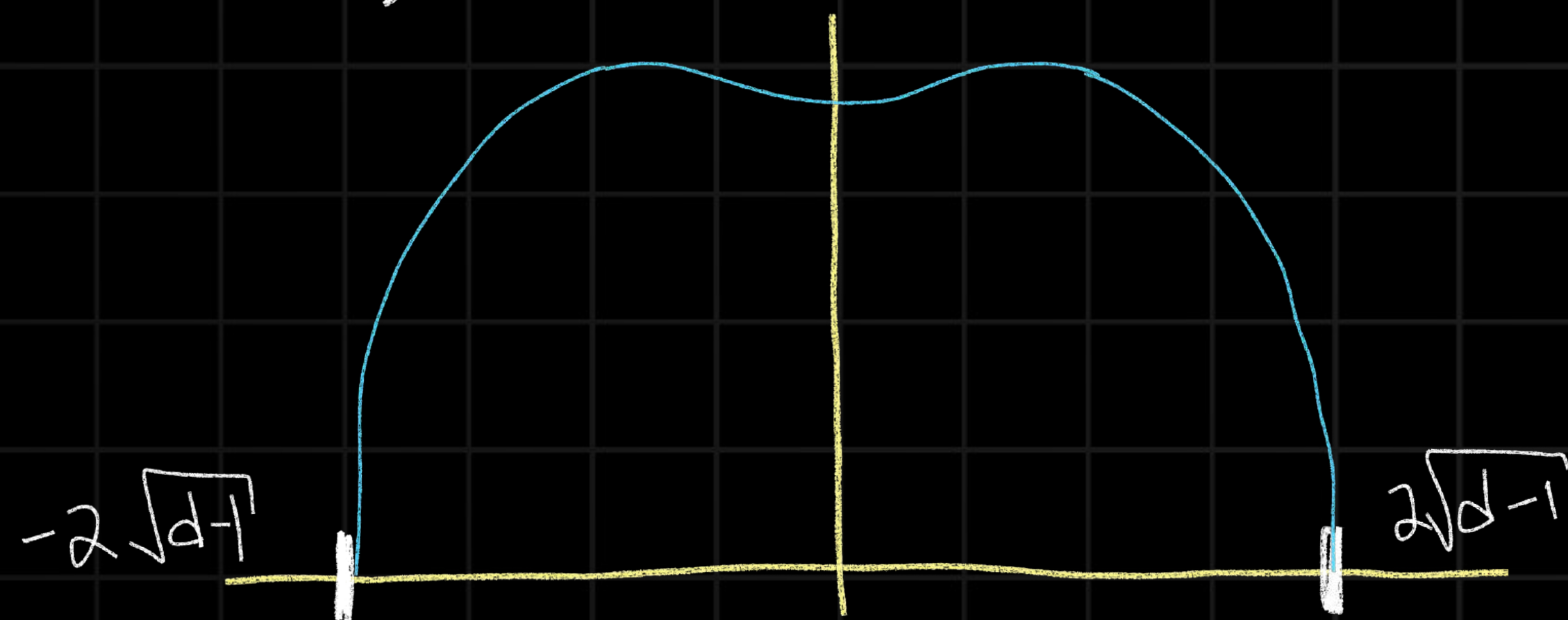
Alon-Boppana: $\lambda(G) \geq 2\sqrt{d-1} - o(1)$

What is $2\sqrt{d-1}$?

- Spectral radius of T_d (infinite d-ary tree)
- Support of KM_d (Kesten-McKay with param d)

$$KM_d = \underbrace{\mu \oplus \mu \oplus \dots \oplus \mu}_{d \text{ times}},$$

$$\mu = \frac{1}{2}(\delta_1 + \delta_{-1})$$



MSS Overview

$$KM_d = \underbrace{\mu \oplus \mu \oplus \dots \oplus \mu}_{d \text{ times}},$$

$$\mu = \frac{1}{2}(\delta_1 + \delta_{-1})$$

What is the graph analog?

MSS Overview

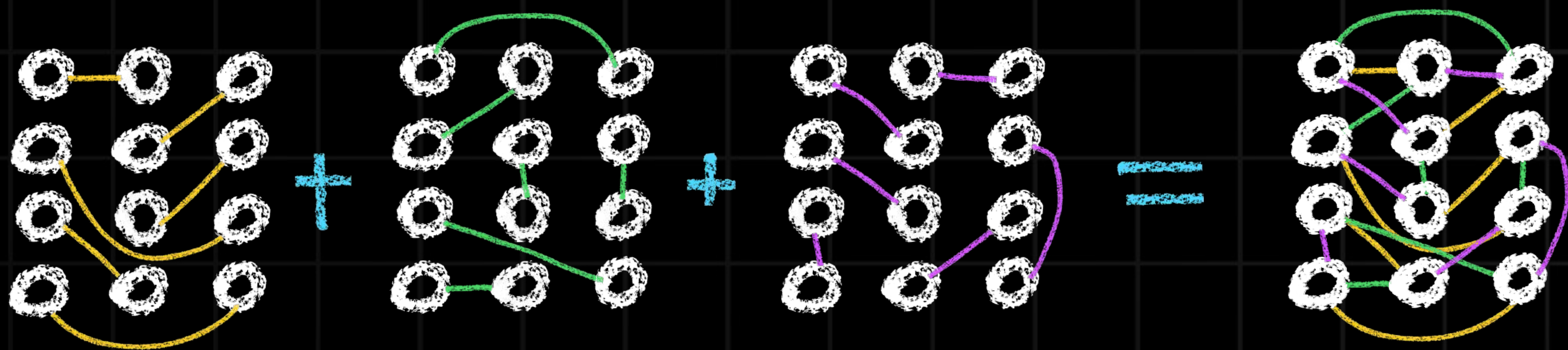
$$KM_d = \mu \boxplus \mu \boxplus \dots \boxplus \mu,$$

$$\mu = \frac{1}{2}(\delta_1 + \delta_{-1})$$

Graph:

$$G_P = M_1 + M_2 + \dots + M_d,$$

$$M_i = P_i^T M P_i$$



Thm: $\lambda_2(G_P) \leq \max\text{-support}(KM_d)$

($\exists P$ s.t.)

(Interlacing & Quadrature)

MSS Inspired Constructions

$$- G_P = M_1 + M_2 + \dots + M_d \quad M_i = P_i^T M P_i$$

$$[MSS'15] \quad \lambda_2(G_P) \leq \text{max-support}(\mu_M^{\boxplus d}) = 2\sqrt{d-1}$$

$$- G_P = G_1 + G_2 + \dots + G_t \quad G_i = P_i^T G P_i$$

$$\lambda_2(G_P) \leq \text{max-support}(\mu_G^{\boxplus t})$$

$$- G_P = G_1 G_2 \dots G_t^2 \dots G_2 G_1 \quad \text{degree } D = d^{2t}$$

$$\lambda_2(G_P) \leq \text{max-support}(\mu_{G^2}^{\boxplus t}) = ?$$

?

MSS Inspired Constructions

$$- G_P = M_1 + M_2 + \dots + M_d \quad M_i = P_i^T M P_i$$

$$[MSS'15] \quad \lambda_2(G_P) \leq \text{max-support}(M_M^{\boxplus d}) = 2\sqrt{d-1}$$

$$- G_P = G_1 + G_2 + \dots + G_t \quad G_i = P_i^T G P_i$$

$$\lambda_2(G_P) \leq \text{max-support}(M_G^{\boxplus t})$$

$$- G_P = G_1 G_2 \dots G_t^2 \dots G_2 G_1$$

$$\lambda_2(G_P) \leq \text{max-support}(M_{G^2}^{\boxplus t}) \approx !$$

Proving I&Q for $\boxtimes$

Reminder: the convolution recipe

$$\mu \longrightarrow G_\mu(x)$$

$$\downarrow$$
$$R_\mu(y)$$

$$\downarrow$$
$$R_{\mu^{\boxplus d}} = d \cdot R_\mu$$

$$\downarrow$$
$$G_{\mu^{\boxplus d}}(x) \longleftarrow \mu^{\boxplus d}$$

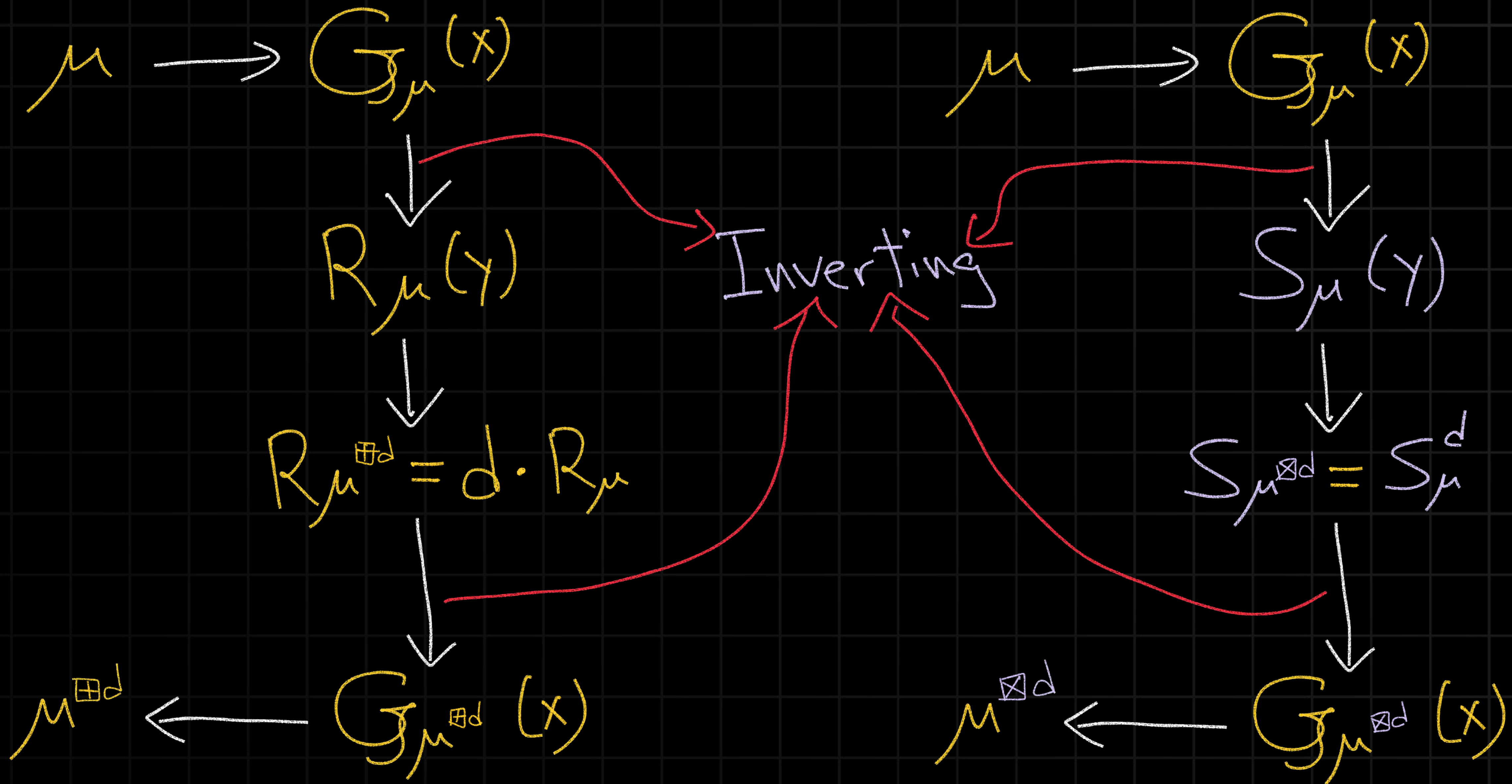
$$\mu \longrightarrow G_\mu(x)$$

$$\downarrow$$
$$S_\mu(y)$$

$$\downarrow$$
$$S_{\mu^{\boxtimes d}} = S_\mu^d$$

$$\downarrow$$
$$G_{\mu^{\boxtimes d}}(x) \longleftarrow \mu^{\boxtimes d}$$

Reminder: the convolution recipe



Reminder: the convolution recipe

$$\mu \longrightarrow G_\mu(x)$$

$$\downarrow$$
$$R_\mu(y)$$

$$\downarrow$$
$$R_{\mu^{\boxplus d}} = d \cdot R_\mu$$

[MSS'15] Bounding $\mu^{\boxplus d}$

$$\mu^{\boxplus d} \longleftarrow G_{\mu^{\boxplus d}}(x)$$

$$\mu \longrightarrow G_\mu(x)$$

$$\downarrow$$
$$S_\mu(y)$$

$$\downarrow$$
$$S_{\mu^{\boxtimes d}} = S_\mu^d$$

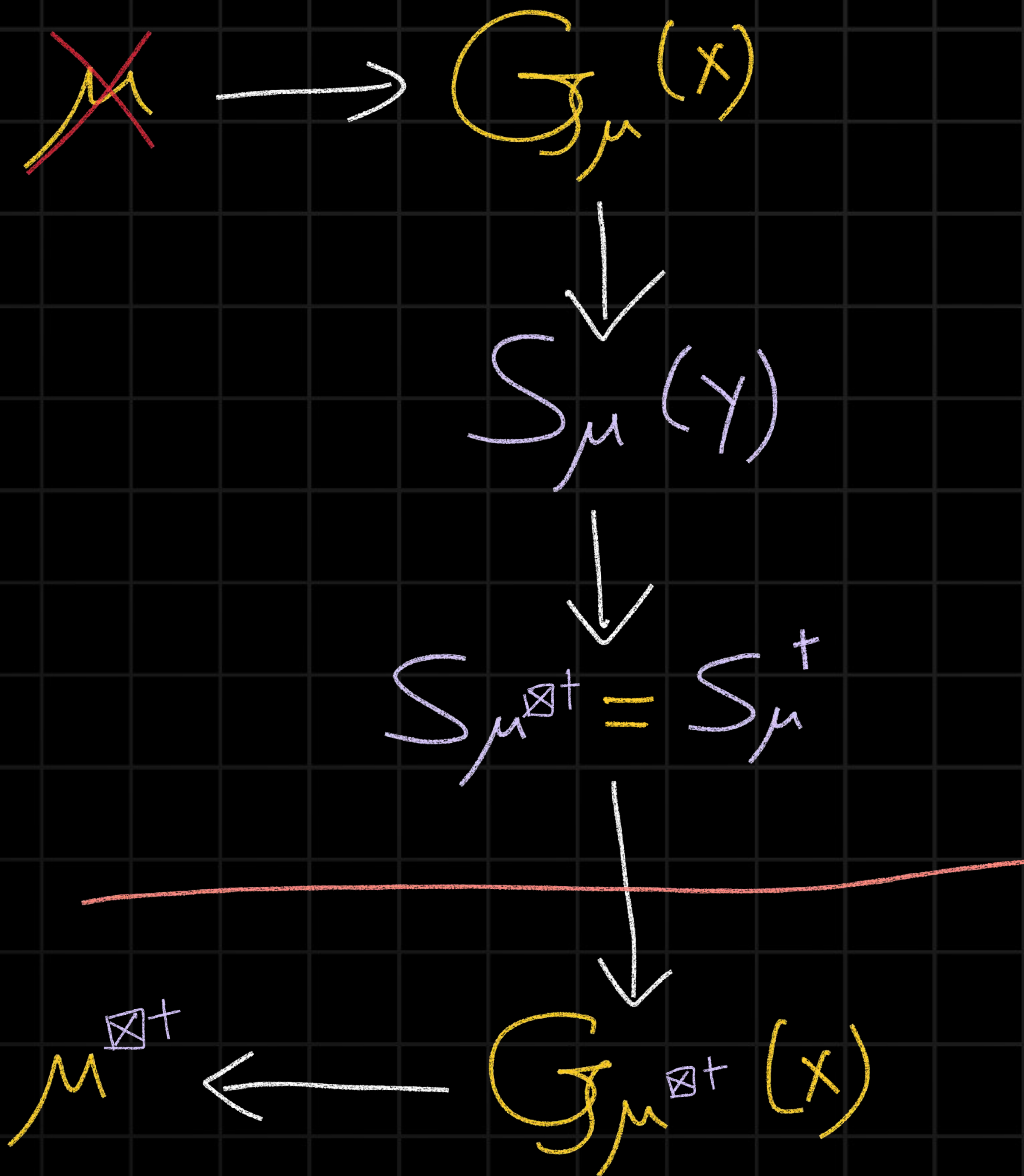
$$\mu^{\boxtimes d} \longleftarrow G_{\mu^{\boxtimes d}}(x)$$

Rotating Expanders

$$G_P = G_1 G_2 \cdots G_t^2 \cdots G_2 G_1$$

Idea (the Adapter)

- Use KM_d instead of μ .
- Prove $S_\mu \approx S_{KM_d}$



The Adapter

Goal: $G_\mu(x) - G_{KM}(x)$ is small.

Ideally, $O_n(1)$
↑

Recall: $G(x) = \sum_{r=0}^{\infty} \frac{m_r}{x^{r+1}}$

For G with eigenvalues μ ,

$$m_r(\mu) = \frac{1}{n} \cdot \#\{\text{closed walks of length } r\}$$

The Adapter: Proof Sketch

Recall: $G(x) = \sum_{r=0}^{\infty} \frac{m_r}{x^{r+1}}$

$$|G_{\mu}(x) - G_{KM}(x)| = \left| \sum_{r=0}^{\infty} \frac{m_r(\mu) - m_r(KM)}{x^{r+1}} \right|$$

$$\leq \underbrace{\left| \sum_{r=0}^h \frac{m_r(\mu) - m_r(KM)}{x^{r+1}} \right|}_{\text{taking } h = \text{girth}} + \underbrace{\left| \sum_{r=h+1}^{\infty} \frac{m_r(\mu) - m_r(KM)}{x^{r+1}} \right|}_{\frac{2}{x-\lambda} \cdot \left(\frac{\lambda}{x}\right)^h}$$

taking $h = \text{girth}$

$$m_r(KM) = \# \text{ in } T_d$$

[McKay'81]

The Adapter: Proof Sketch

Recall: $G(x) = \sum_{r=0}^{\infty} \frac{m_r}{x^{r+1}}$

$$G_{\mu}(x) - G_{KM}(x) = \sum_{r=0}^{\infty} \frac{m_r(\mu) - m_r(KM)}{x^{r+1}}$$

$$\leq \underbrace{\left| \sum_{r=0}^h \frac{m_r(\mu) - m_r(KM)}{x^{r+1}} \right|}_{\text{taking } h = \text{girth}} + \underbrace{\left| \sum_{r=h+1}^{\infty} \frac{m_r(\mu) - m_r(KM)}{x^{r+1}} \right|}_{\frac{2}{x-\lambda} \cdot \left(\frac{\lambda}{x}\right)^h}$$

taking $h = \text{girth}$

$$m_r(KM) = \# \text{ in } T_d$$

[McKay'81]

$$\frac{2}{x-\lambda} \cdot \left(\frac{\lambda}{x}\right)^h$$

depends on G !

The Adapter: Proof Sketch

Recall: $G(x) = \sum_{r=0}^{\infty} \frac{m_r}{x^{r+1}}$

$$G_{\mu}(x) - G_{KM}(x) = \sum_{r=0}^{\infty} \frac{m_r(\mu) - m_r(KM)}{x^{r+1}}$$

$$\leq \underbrace{\left| \sum_{r=0}^h \frac{m_r(\mu) - m_r(KM)}{x^{r+1}} \right|}_{0} + \underbrace{\left| \sum_{r=h+1}^{\infty} \frac{m_r(\mu) - m_r(KM)}{x^{r+1}} \right|}_{\frac{2}{x-\lambda} \cdot \left(\frac{\lambda}{x}\right)^h}$$

taking $h = \text{girth}$

$$m_r(KM) = \# \text{ in } T_d$$

[McKay'81]

(getting S_{μ} is calculus)

MSS Inspired Construction I Rotating Expanders

$$- G_P = G_1 G_2 \cdots G_t^2 \cdots G_2 G_1, \quad G_i = P_i^T G P_i, \quad D = d^{2t}$$

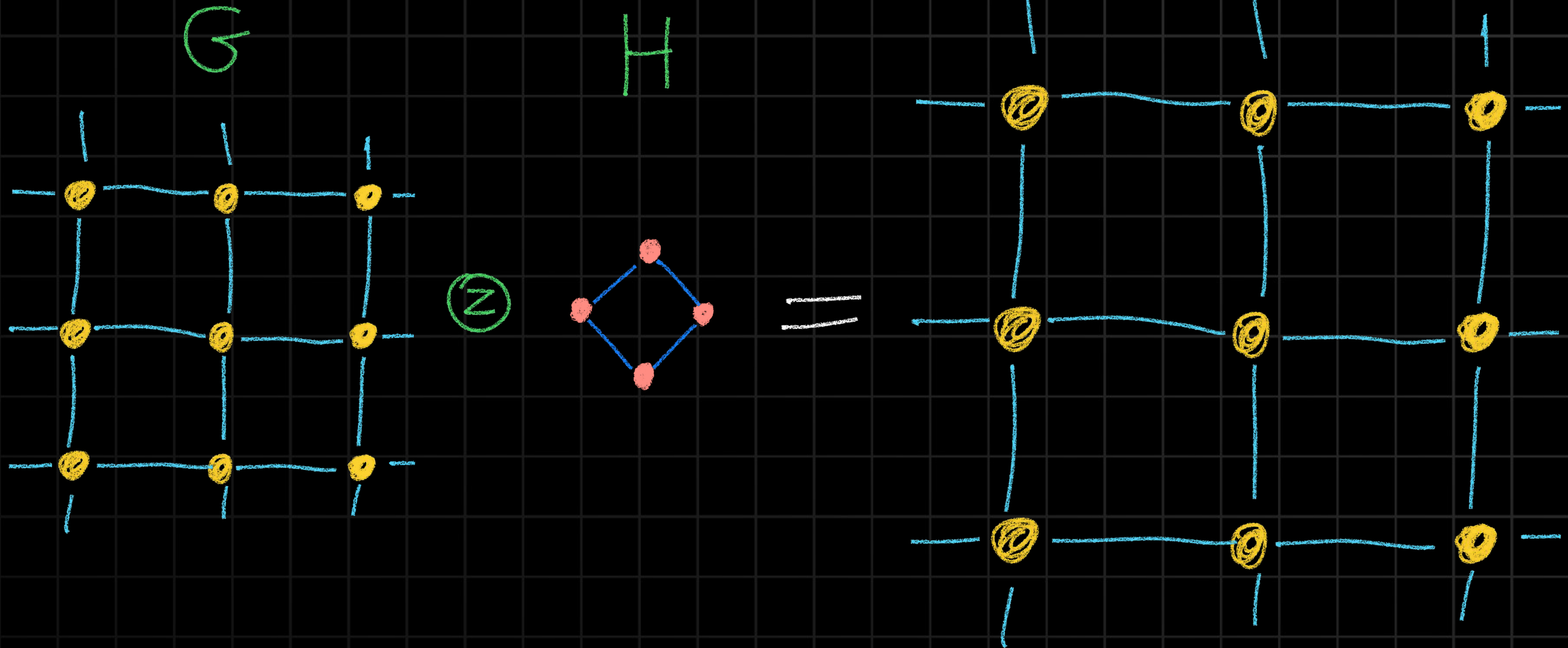
Thm [CM'23]: for G Ramanujan,

$$\lambda(G_P) \leq \underbrace{\frac{(t+1)^{t+1}}{t^t}}_{\text{Fuss-Catalan growth rate}} d \cdot (d-1)^{t-1} + O(t \cdot 2^{-h}) < \underbrace{e \cdot (t+1)}_{\ll 2^t} \sqrt{D}$$

($\exists P$ s.t.)

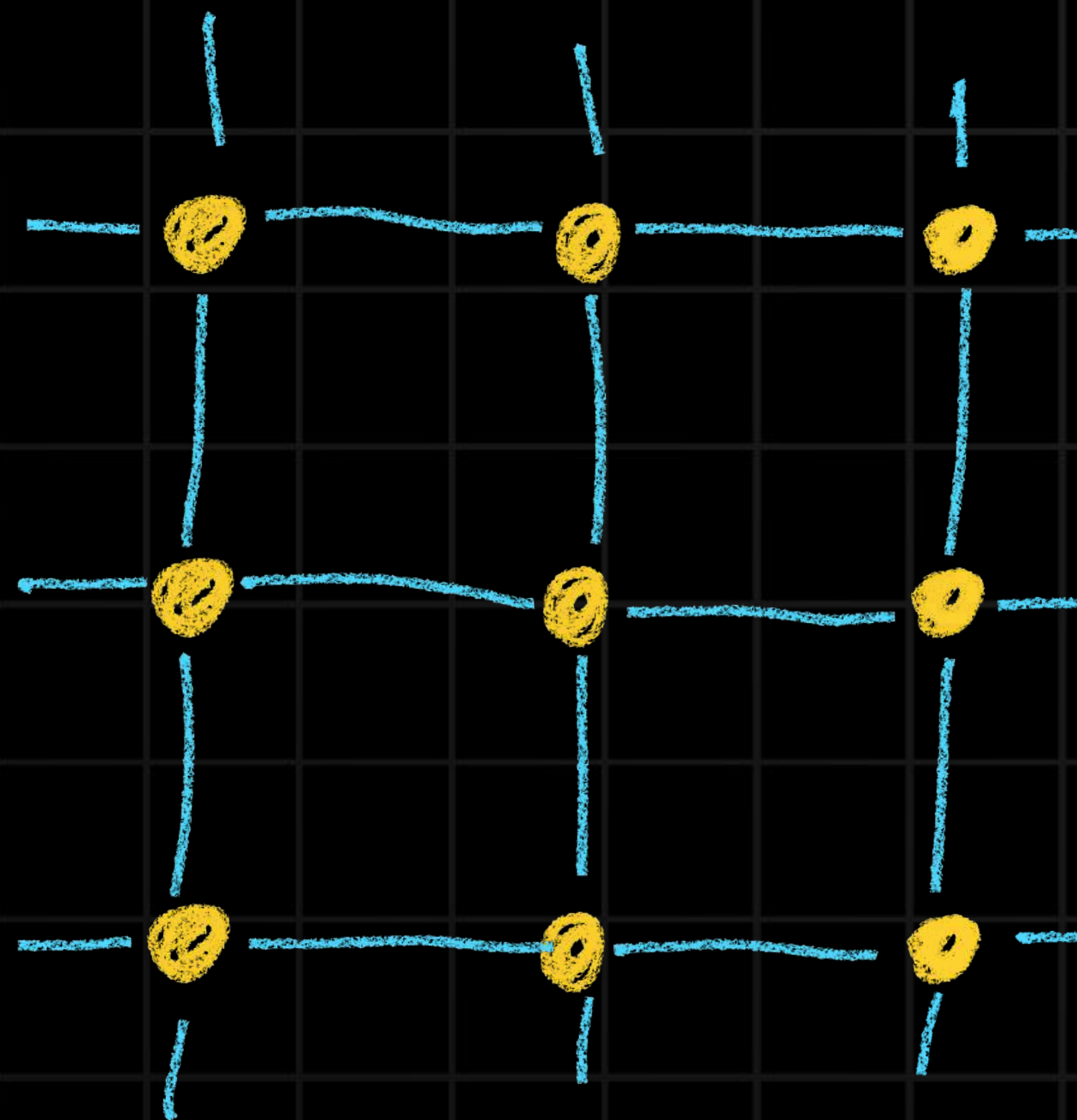
Tight! Lower bound proven in [CCMP'24]

The Zig-Zag product $G \circledast H$

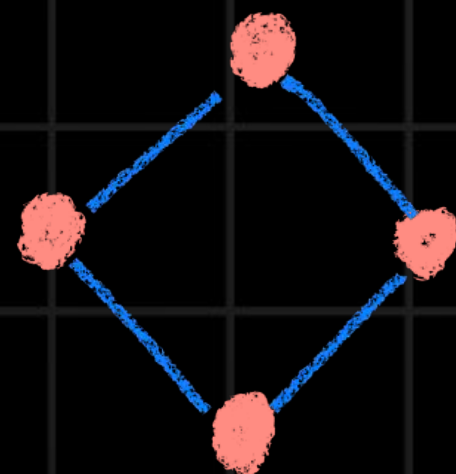


The Zig-Zag product $G \circledast H$

G



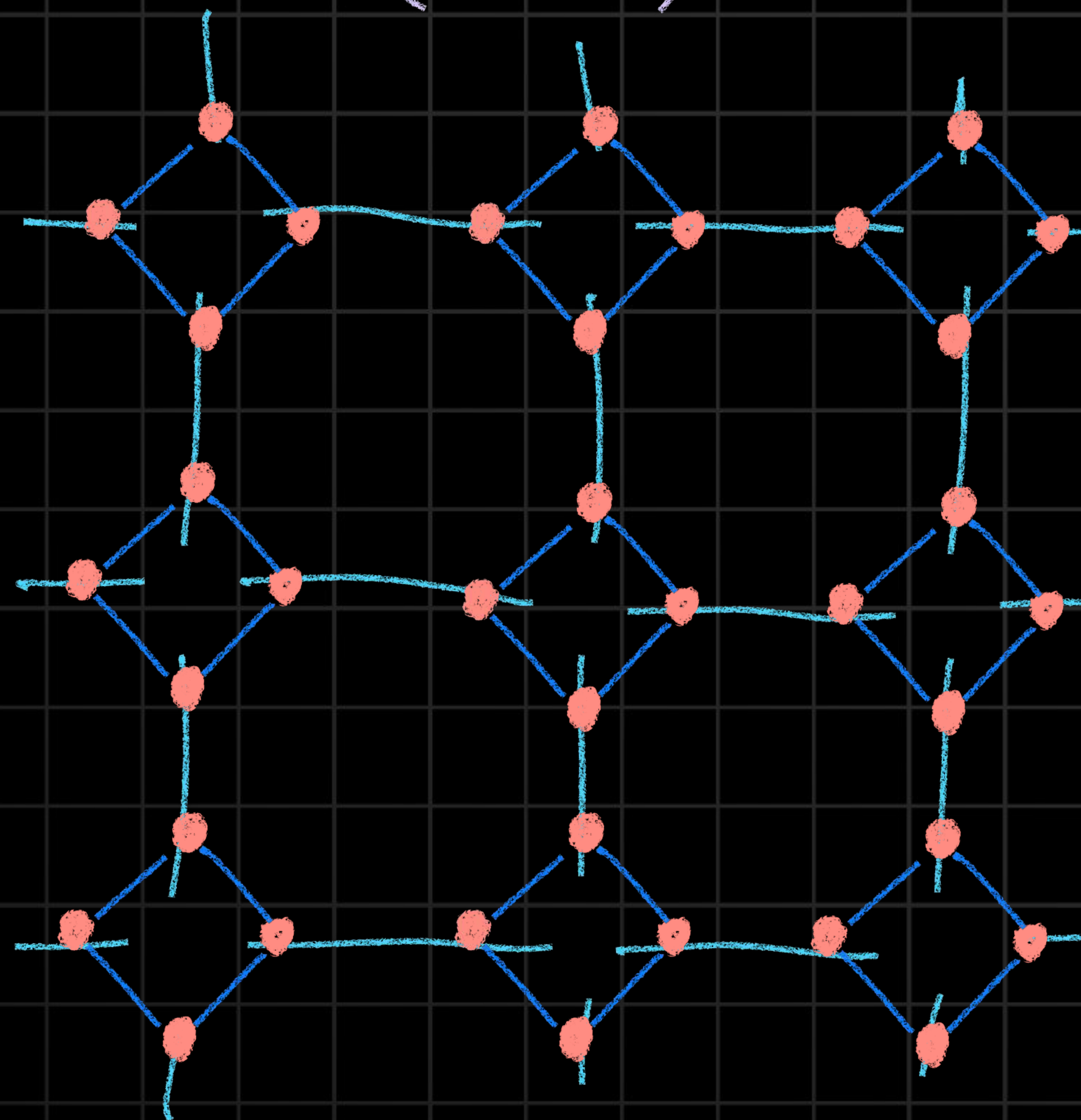
H



$\circledast$

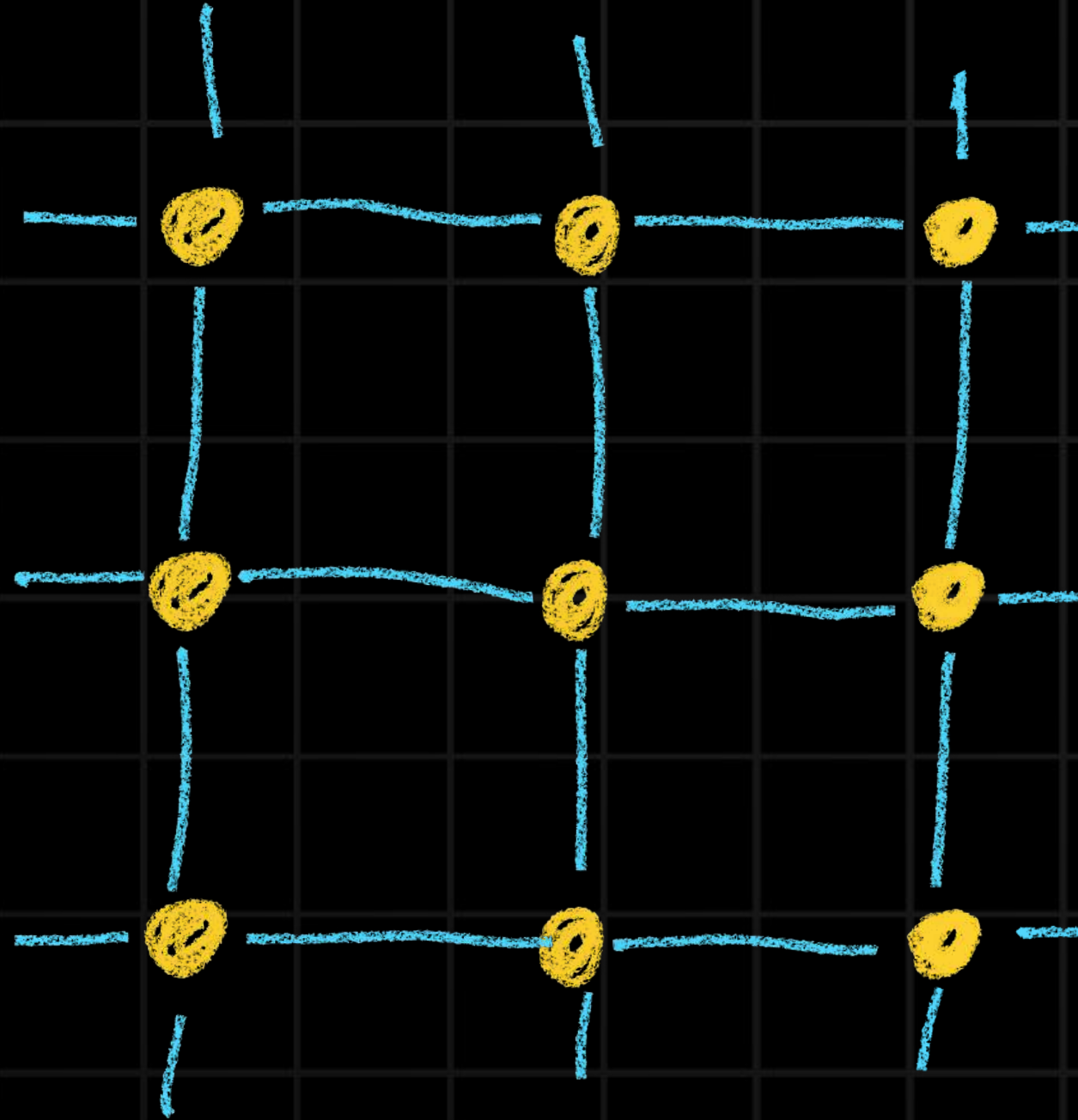
$=$

$(G \circledast H)$

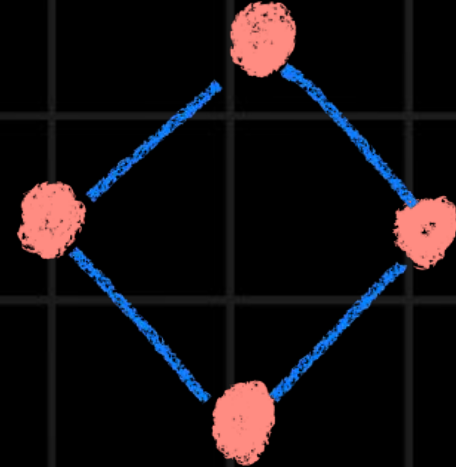


The Zig-Zag product $G \circledast H$

G

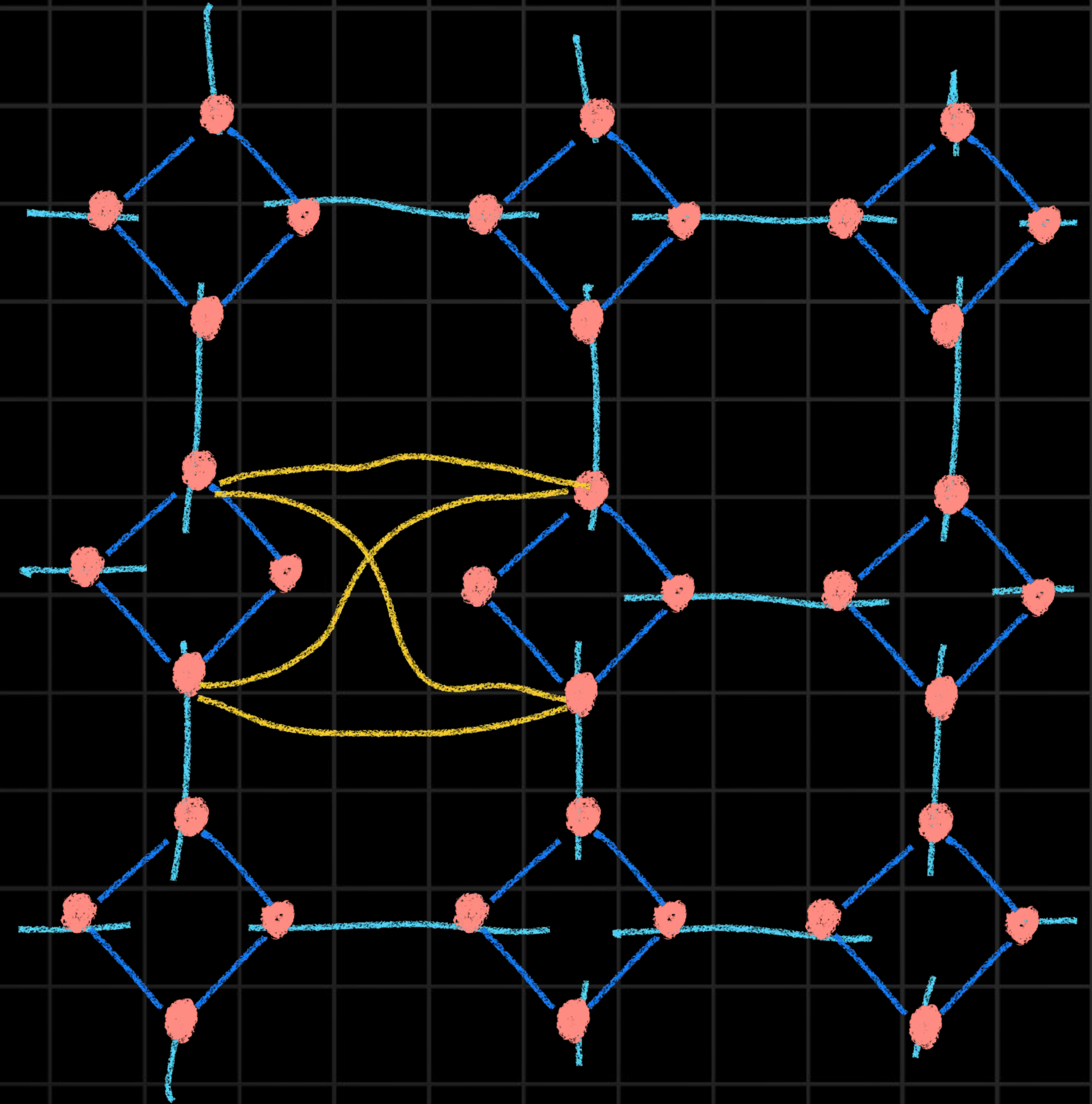


H



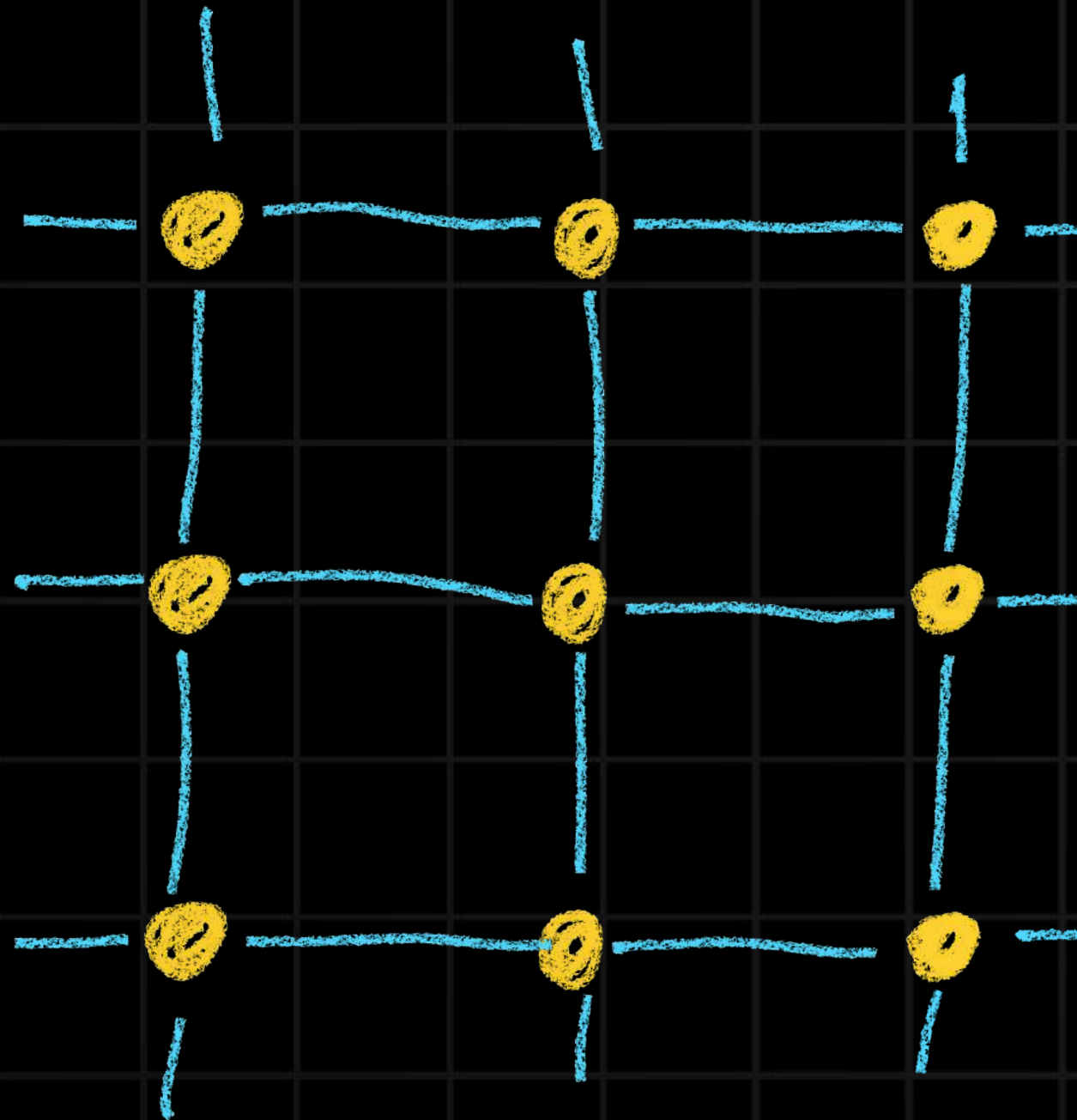
$\circledast$

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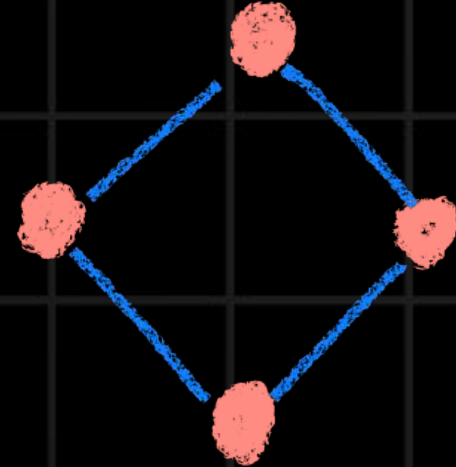


The Zig-Zag product $G \circledast H$

G

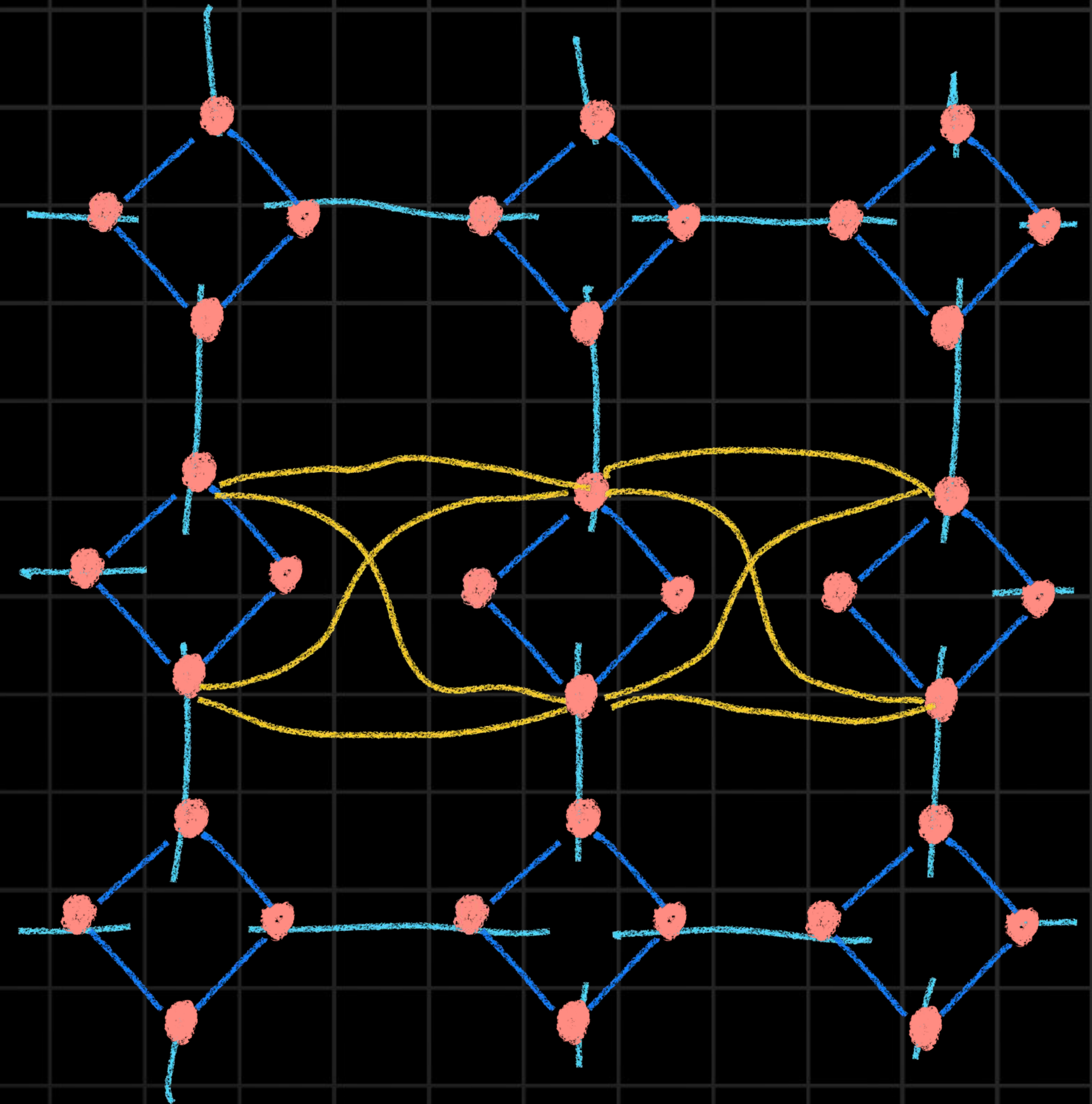


H



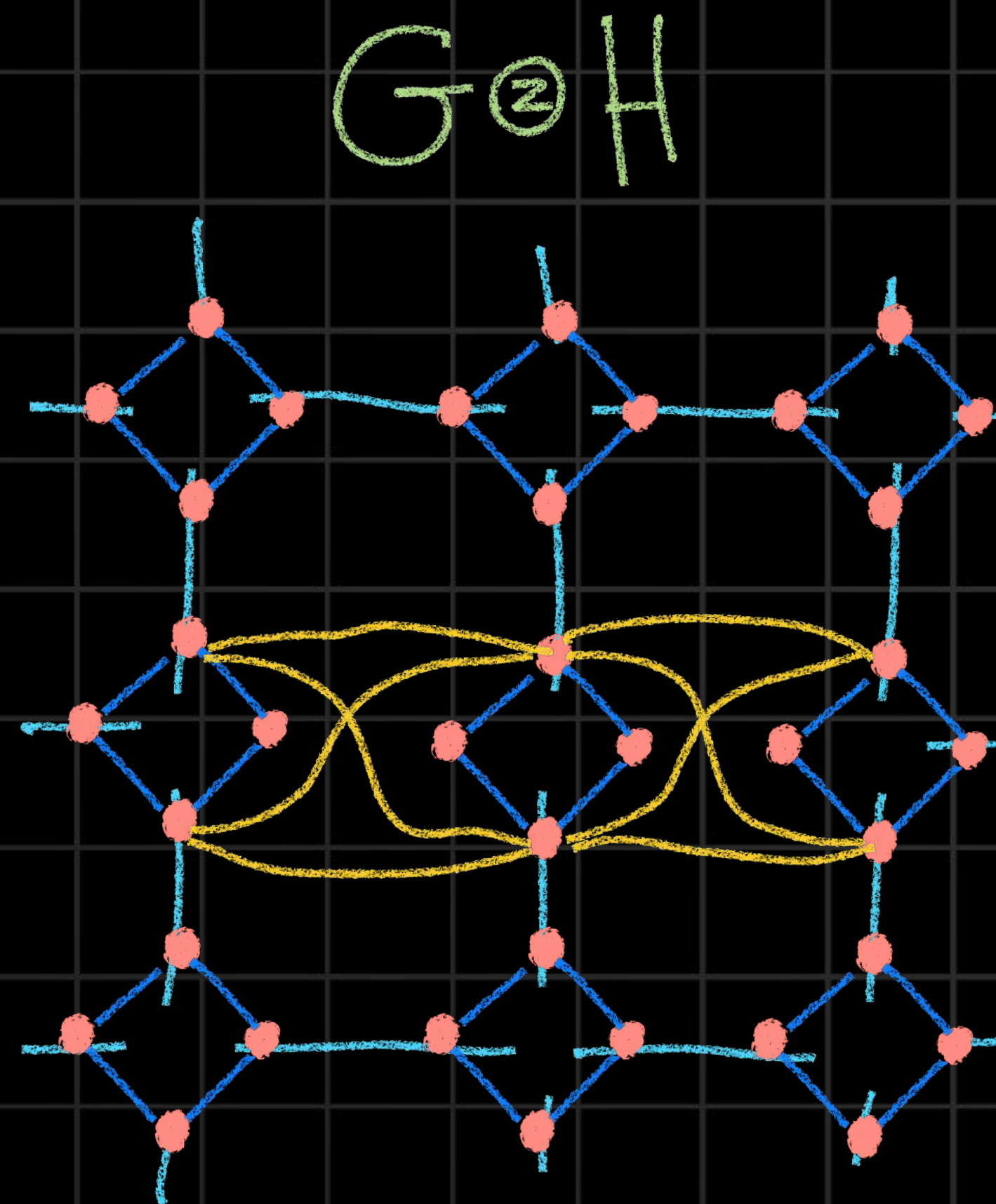
$\circledast$

$=$



Properties of ②

	Size	D	ω
G	n	d	ω_G
H	d	c	ω_H
$G \circledast H$	nd	c^2	$\leq \omega_G + \omega_H$ [RVW] $= O(D^{-1/4})$



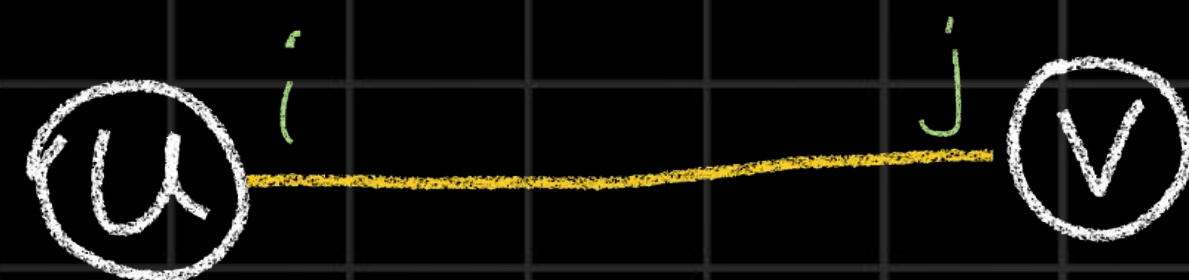
Analyzing $G \otimes H$ in Free Probability

$$G \otimes H = \tilde{H} \dot{G} \tilde{H}$$

$$\tilde{H} = I_n \otimes H = \begin{pmatrix} H & & \\ & H & \\ & & \ddots \\ & & & H \end{pmatrix}$$

Random G : $\tilde{H} P^T M P \tilde{H}$

arbitrary
matching, ± 1



Thm [CCM'24]:

$$\lambda_2(G \otimes H) \leq \max\text{-support}(\mu_{H^2} \boxtimes \mu_M)$$

($\exists G$ s.t.)

$$\dot{G}_{(u,i),(v,j)} = 1$$

Analyzing $G \boxplus H$ in Free Probability

Thm [CCM'24]: $\lambda_2(G \boxplus H) \leq \underbrace{\text{max-support}(\mu_{H^2} \boxtimes \mu_M)}_{Z_H}$
($\exists G$ s.t.)

Moreover, this is a lower bound:

$$\lambda(G \boxplus H) \geq Z_H - o_n(1)$$

Proven using
analytic combinatorics.

Analyzing $G \boxplus H$ in Free Probability

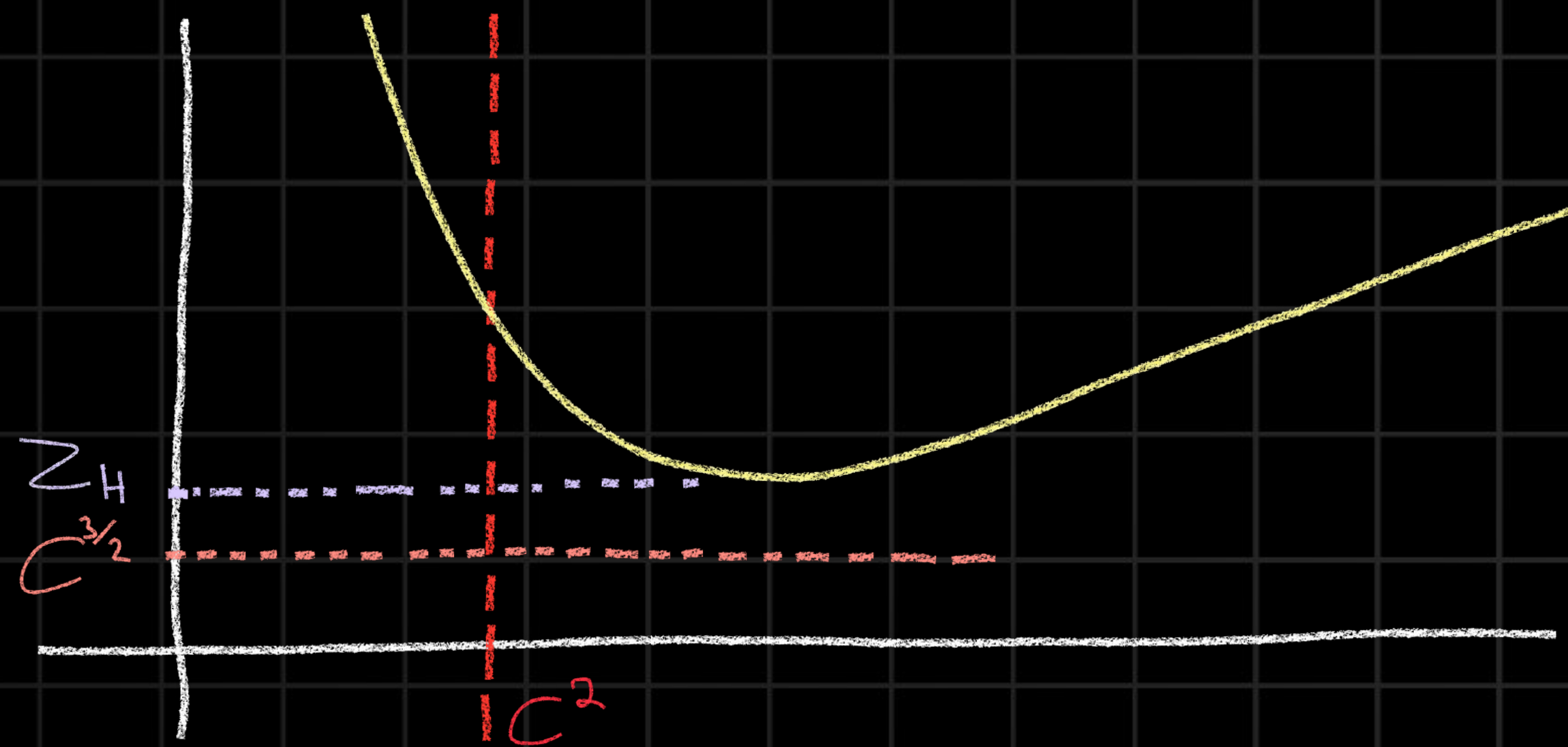
Thm [CCM'24]: $\lambda_2(G \boxplus H) \leq \underbrace{\max\text{-support}}_{Z_H}(\mu_{H^2} \boxtimes \mu_M)$
($\exists G$ s.t.)

Moreover, this is a lower bound:

$$\lambda(G \boxplus H) \geq Z_H - O_n(1)$$

$$\omega_{G \boxplus H} \geq \frac{C^{3/2}}{C^2} = C^{-1/2} = D^{-1/4}$$

$$Z_H = \min_{x > C^2} x \sqrt{1 - \frac{1}{x \cdot G_{H^2}(x)}}$$



Analyzing $G \otimes H$ in Free Probability

Thm [CCM'24]: $\lambda_2(G \otimes H) \leq \underbrace{\text{max-support}(\mu_{H^2} \boxtimes \mu_M)}_{Z_H}$
($\exists G$ s.t.)

Moreover, this is a lower bound:

$$\lambda(G \otimes H) \geq Z_H - o_n(1)$$

Corollary:

$$w(G \otimes H) \geq D^{-1/4}$$

(and not $D^{-1/2}$)

The [RVW] bound
is asymptotically tight

Recap

- Using FP, proved tight bounds for the ZigZag product $\textcircled{z}$ and others (derand. square $\textcircled{s}$, replacement product $\textcircled{r}$).
- "Free Graphs" yield better expansion properties (rotating expanders)

Questions?

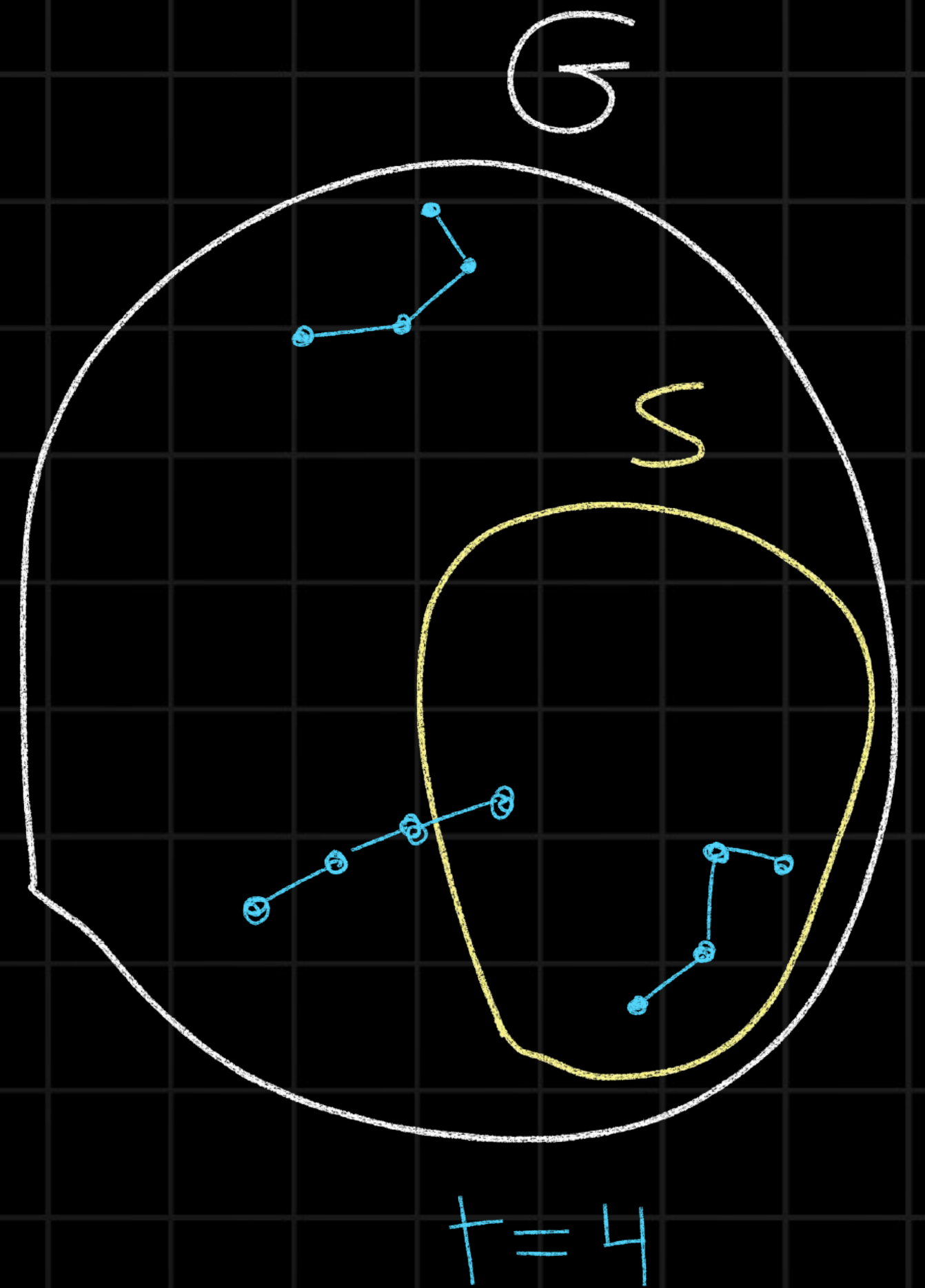
Open Problem I

Trapped Random Walks

(or: the Expander Hitting Property)

G is an w -expander. $|S| = \mu \cdot |V|$

$$\Pr[\text{walk} \subseteq S] = ?$$



Open Problem I

Trapped Random Walks

(or: the Expander Hitting Property)

G is an w -expander. $|S| = \mu \cdot |V|$

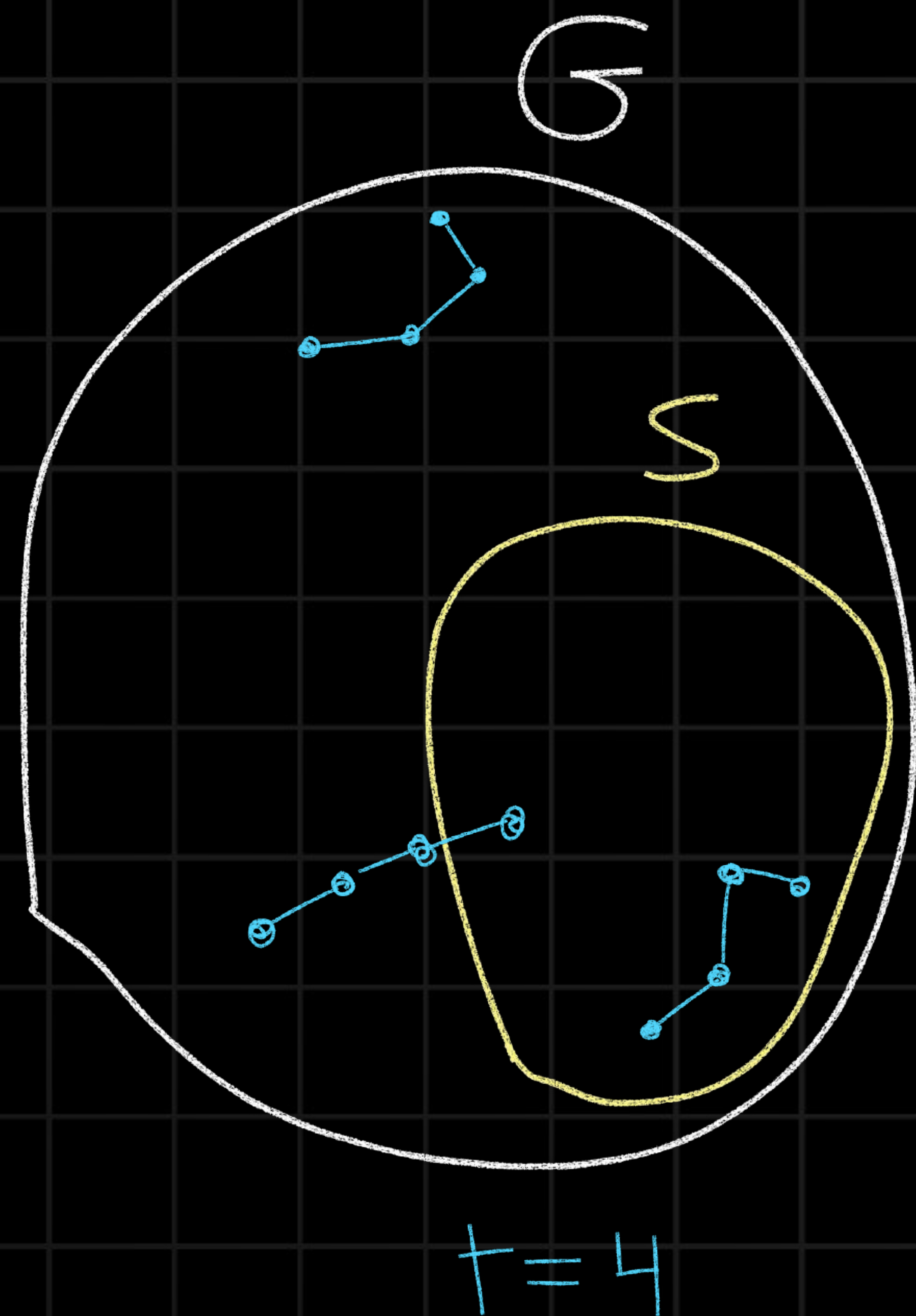
$$\Pr[\text{path} \subseteq S] \leq \mu \cdot (\mu + w)^{t-1}$$

(Ideal: μ^t)

- The ideal requires $t \cdot \log n$ bits.
- Random Walk: $\log n + (t-1) \cdot \log d$.

Taking $w = \mu$:
($d \approx 4/\mu^2$)

$$\Pr[\text{path} \subseteq S] \approx 2^t \mu^t$$



Open Problem I

Trapped Random Walks

(or: the Expander Hitting Property)

$$\Pr[\text{walk} \subseteq S] \leq \mu \cdot (\mu + \omega)^{t-1}$$

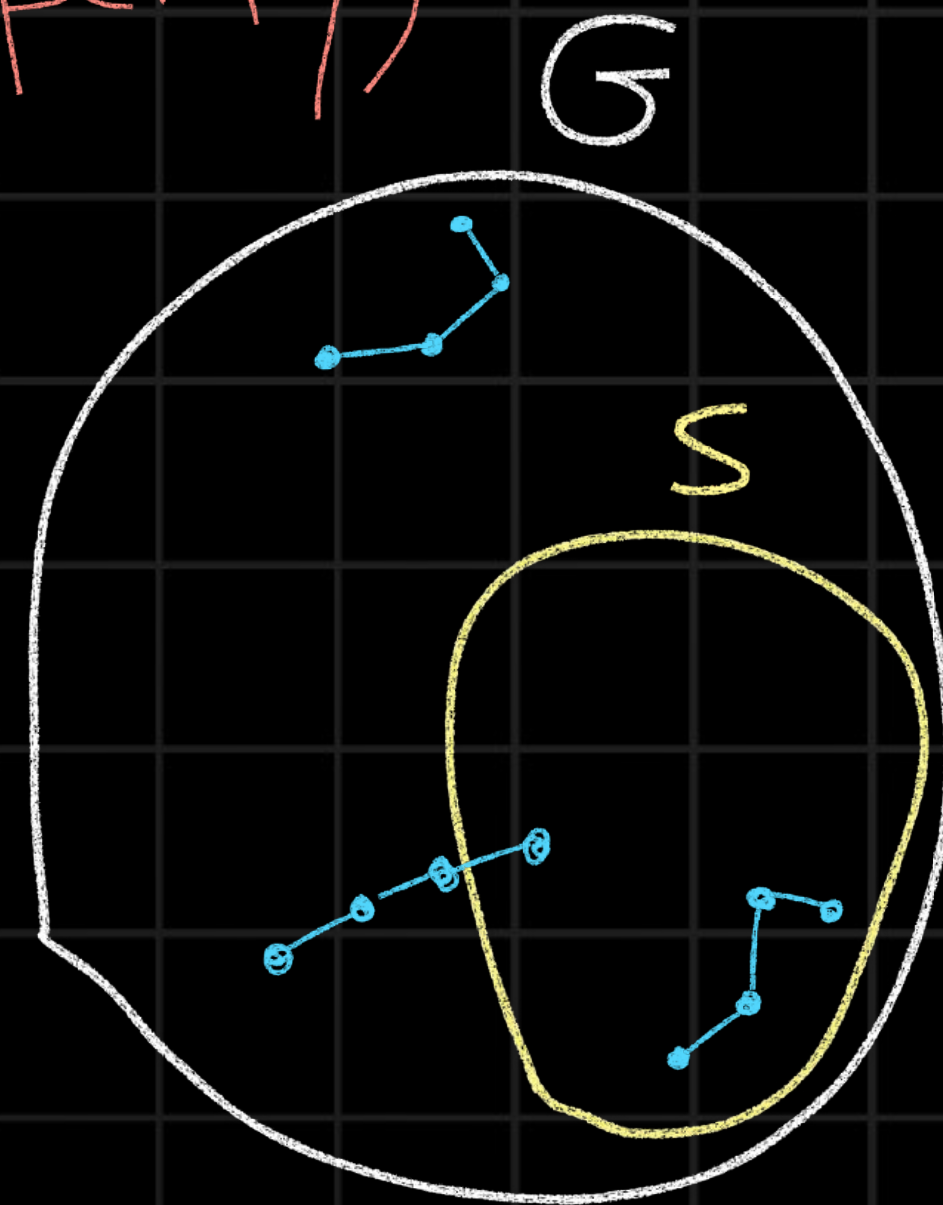
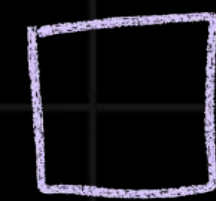
Proof: $\frac{1}{n} \mathbf{1}^T S G S G S G S \mathbf{1}$ ($S_{ii} = 1$ for $i \in S$)

$$= \frac{1}{n} \mathbf{1}^T S^2 G S^2 G S^2 G S^2 \mathbf{1}$$

$$\leq \mu \cdot \|S G S\| \cdot \|S G S\| \cdot \|S G S\|$$

$$\left(\frac{1}{n} \mathbf{1}^T S^2 \mathbf{1} = \mu \right)$$

$$\|S G S\| \leq \underbrace{\|S J S\|}_{\mu} + \underbrace{\|S E S\|}_{\leq \omega}$$



Open Problem I Hitting Property for Rotating Expanders

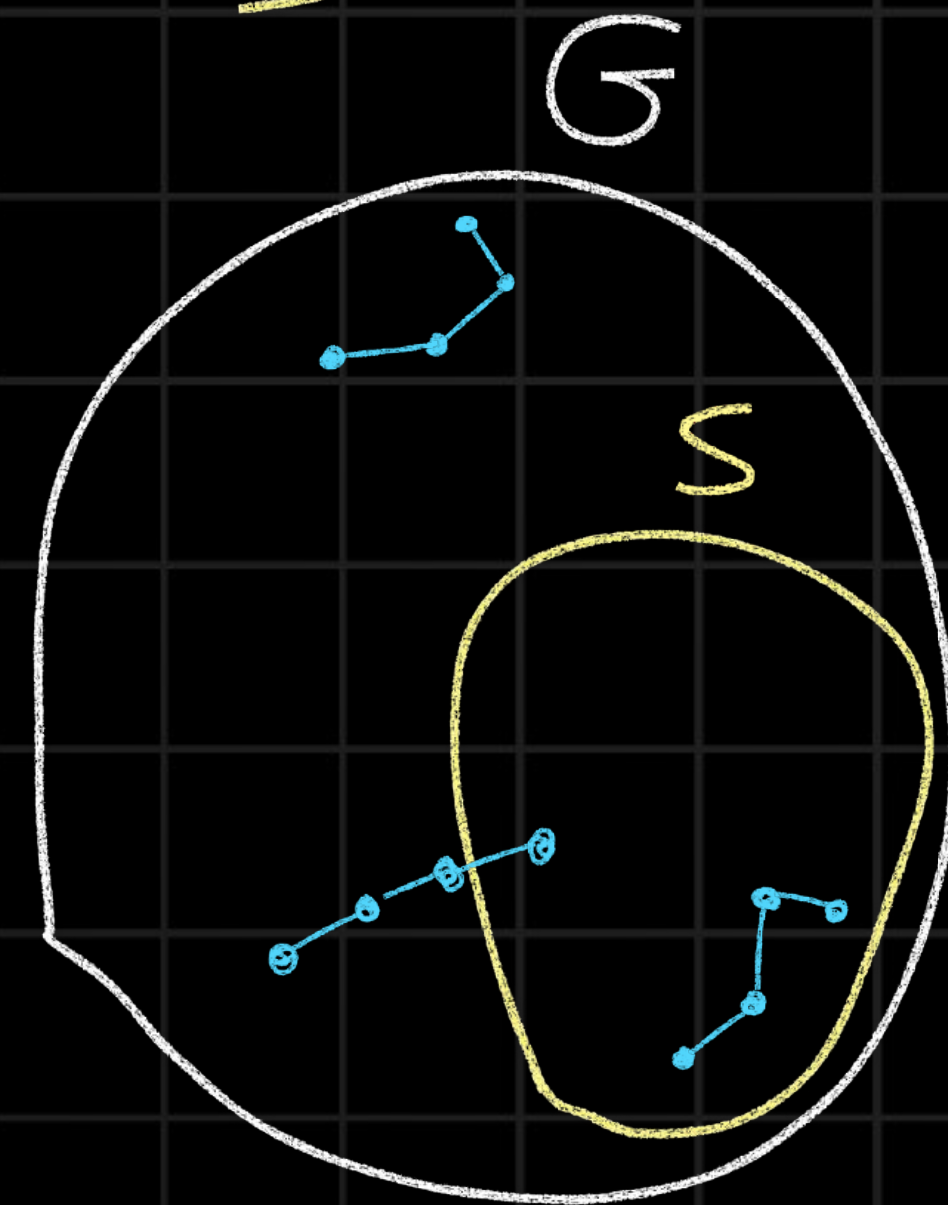
$$\Pr \left[\begin{array}{c} \text{path} \\ \parallel \\ \end{array} \subseteq S \right] \leq \mu \cdot (\mu + w)^{t-1}$$

Proof: $\frac{1}{n} \mathbf{1}^T S G_3 S G_2 S G_1 S \mathbf{1}$

$$= \frac{1}{n} \mathbf{1}^T S^2 G_3 S^2 G_2 S^2 G_1 S^2 \mathbf{1}$$

$$\leq \mu \cdot \|S G_3 S\| \cdot \|S G_2 S\| \cdot \|S G_1 S\|$$

$$\|S G_i S\| \leq \underbrace{\|S J S\|}_{\mu} + \underbrace{\|S E_i S\|}_{\leq ?}$$



Open Problem I Hitting Property for Rotating Expanders

$$\Pr[\text{path} \subseteq S] \leq \mu \cdot (\mu + \omega)^{t-1}$$

Proof: $\frac{1}{n} \mathbf{1}^T S G_3 S G_2 S G_1 S \mathbf{1}$

$$= \frac{1}{n} \mathbf{1}^T S^2 G_3 S^2 G_2 S^2 G_1 S^2 \mathbf{1}$$

or, analyze this term.

$$\leq \mu \cdot \|S G_3 S\| \cdot \|S G_2 S\| \cdot \|S G_1 S\|$$

for $t=3$: $\mu(\mu + \frac{\omega}{\sqrt{2}})^2$

$$\|S G_i S\| \leq \|S J S\| + \|S E_i S\|$$

We suspect $\mu(\mu + \frac{\omega}{\sqrt{t-1}})^{t-1}$

Open Problem II

Small-Bias sets

Goal: find a **small** set $S \subseteq \{\pm 1\}^n$ that ε -fools all linear tests.

$$J \subseteq [n] \quad f_J(x) = \prod_{j \in J} x_j \quad x \in \{\pm 1\}^n$$

$$\forall J \quad \mathbb{E}_{x \in \{\pm 1\}^n} f_J(x) = 0$$

$$|S| = ?$$

$$\left| \mathbb{E}_{x \in S} f_J(x) \right| \leq \varepsilon$$

Open Problem II

Small-Bias sets

Goal: find a **small** set $S \subseteq \{\pm 1\}^n$ that ε -fools all linear tests.

$$J \subseteq [n] \quad f_J(x) = \prod_{j \in J} x_j$$

$$|S| = \Omega\left(\frac{n}{\varepsilon^2 \cdot \log \frac{1}{\varepsilon}}\right)$$

$$\forall J \quad \mathbb{E}_{x \in \{\pm 1\}^n} f_J(x) = 0$$

non-explicit:

$$\left| \mathbb{E}_{x \in S} f_J(x) \right| \leq \varepsilon$$

$$|S| = O\left(\frac{n}{\varepsilon^2}\right)$$

Open Problem II

Small-Bias sets

Amplification using expanders: Let $|S| = m$, ϵ_0 -biased.

G a d -regular w -expander. New set S_+ :

$$X_1 - X_2 - \dots - X_t \in G$$

$$|S_+| = m \cdot d^{t-1}$$

$$y \in S_+ : y_i = x_{1i} \cdot x_{2i} \cdots x_{ti}$$

$J \subseteq [n]$, P diag. matrix, $P_{xx} = f_J(x)$.

$$\epsilon_+ = \frac{1}{n} \mathbf{1}^T \underbrace{P G P G P G P}_t \mathbf{1} \leq O\left((\epsilon_0^2 + w)^{t/2}\right)$$

Open Problem II

Small-Bias sets

$$|S| = m, \varepsilon_0\text{-biased}, |S_t| = m \cdot d^{t-1}, \varepsilon_t \leq O\left((\varepsilon_0^2 + w)^{t/2}\right)$$

Enables constructing ε -biased T of size $O\left(\frac{n}{\varepsilon^4}\right)$

What if we analyzed

$$\varepsilon_t = \frac{1}{n} \mathbf{1}^T P G_3 P G_2 P G_1 P \mathbf{1} \quad ?$$

Thank You

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